Syntax-Guided Program Synthesis

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Syntax-Guided Search-based Program Synthesis



Motivating Applications

- Superoptimizing compilers: Given a program fragment P, find a functionally equivalent program with resource constraints (e.g. fewer instructions, or avoid certain expensive instructions)
- Program repair: Automatically edit a program locally to fix a bug (particularly helpful to students in Intro Programming courses)
- Proof objects for verification: template-guided synthesis of inductive invariants, ranking functions, program analysis rules, ...
- Programming by examples / demonstration: Can non-programmers communicate intent intuitively?

Syntax-Guided Synthesis (SyGuS)



- □ Fix a background theory T: fixes types and operations
- Function to be synthesized: name f along with its type
 - General case: multiple functions to be synthesized
- Inputs to SyGuS problem:
 - Specification \u03c6(x, f(x))
 Typed formula using symbols in T + symbol f
 - Set E of expressions given by a context-free grammar
 Set of candidate expressions that use symbols in T
- Computational problem:
 Output e in E such that \u03c6[f/e] is valid (in theory T)
- Syntax-guided synthesis; FMCAD'13 with Bodik, Juniwal, Martin, Raghothaman, Seshia, Singh, Solar-Lezama, Torlak, Udupa

SyGuS Competition



SyGuS Progress



Over 2800 benchmarks

- Hacker's delight
- Invariant generation (based on verification competition SV-Comp)
- FlashFill (programming by examples system from Microsoft)
- Synthesis of attack-resilient crypto circuits
- Program repair
- Motion planning
- ICFP programming competition
- Special tracks for competition
 - Invariant generation
 - Programming by examples
 - Conditional linear arithmetic
- Current winner: CVC4 (Reynolds et al) search integrated in constraint solving

www.sygus.org

Search and Verify



Concept class: Set E of expressions

Examples: Concrete input values

Counterexample-guided Inductive Synthesis (CEGIS)

Goal: Find f such that for all x in D, $\phi(x, f)$ holds

```
I = { }; /* Interesting set of inputs */
Repeat
Learn: Find f such that for all x in I, φ(f, x) holds
Verify: Check if for all x in D, φ(f, x) holds
If so, return f
If not, find x such that ~ φ(f, x) holds, and add x to I
```

Implementing Search

Given: Specification $\phi(x, f(x))$ Grammar for set E of candidate implementations Finite set I of inputs Find an expression e(x) in E s.t. $\phi(x, e(x))$ holds for all x in I

- Enumerative search with lots of optimizations for pruning
- Symbolic constraints over variables encoding desired expression tree
- Stochastic search in spirit of genetic programming
- Divide and conquer strategies to build sub-expressions
- Partial evaluation to rule candidates before fully expanding them
- Establishing unrealizability of synthesis
- Type-directed enumeration

Acceleration Using Learned Probabilistic Models

- Can we bias the search towards likely programs?
- Step 1: Mine existing solutions to convert given grammar into a probabilistic higher-order grammar
 - Weighted production rules
 - Conditioned on parent and sibling context
 - Transfer learning used to avoid overfitting
- Step 2: Enumerative search to generate expressions in decreasing likelihood
 - Use A* with cost estimation heuristic
 - Integrated with previous optimizations (equivalence-based pruning...)

With W. Lee, K. Heo, and M. Naik (PLDI 2018)

Future Directions

- Beyond SMT Solvers: SyGuS-like back-end focused on efficient search, but decoupled from SMT solvers so as to allow interface with alternative testing / verification tools
- More theories, benchmarks, and applications: tables and relational queries, floating point arithmetic
- Quantitative synthesis and optimization
- Applications in scientific computing ??